# **Reliable Error free Coding for Resource Constraint Soc Design**

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**Abstract:** With the extent of video coding applications in various domains, the need for faster and reliable modelling of video codec's is in raise. The conventional modelling of video codec's were developed for high resource compatibility. As the applications of video coding has moved to low resource devices, a need for the proposal of new coding, streaming and buffering logic in resource constraint environment is required. In this paper an approach for the modelling of video coding process using language of hardware description. The operational functionality of the proposed system is validated over Xilinx targeted FPGA device. A comparative analysis of the developed approach for 2 different search approaches for motion estimation is evaluated over the developed system. The comparative result observations were observed to be improved in estimation accuracy and minimization in processing time.

Keywords: Video codec; coding; steaming; buffering unit; codec modelling.

# I. Introduction

Video processors are developing for various needs of applications. Video has become an essential part for multimedia applications. There are, however, many contradicting limitations in video codec implementations. One challenge is the rapid evolution of compression standards with several different algorithms [1]. This requires programmability that is easy to achieve with processor-based platforms. However, achieving the best power, energy, and silicon area efficiency requires custom hardware implementations [14]. On the other hand, hardware (HW) design is more demanding than software (SW) development, and modifications are very expensive and time consuming [8]-[13]. For example, nonrecurring engineering costs, increase rapidly with each technology generation making frequent HW upgrades less favourable. Software only implementation solves the flexibility and upgradeability problem but is not an optimal solution from a performance versus real time need point of view [15], [16]. In video processing the modelling of SoC unit for multi rate service compatibility is still an open challenge. Various approaches were made in regard to obtain scalable coding approach in video coding [2]-[7]. However in such coding approach the processing error is not concentrated. To improvise the visual quality with demanded scalable rate retaining error limit in this paper a scalable error coding approach is proposed. A digital modelling for coding at different level of service is designed to provide maximum visualization quality to decoded video. To obtain the objective of reliable video coding with variable resource feasibility a new coding structure for video coding in SoC model is presented. To outline the suggested approach in this paper a video coding system designed for modelling of SoC is outlined in section 2. Section 3 outlines the proposed scalable coding approach for data streaming. For the presentation of image projection to higher grid interpolation a resolution projection approach is outlined in section 4. The experimental result for the developed system is outlined in section 5. Section 6 presents the conclusion made for the developed approach.

# II. Video Coding: System Outline

For the designing modelling of proposed approach in a generalized system model is developed. The proposed design architecture for the suggested unit is as shown in below figure.





The system reads the input data from the source file. The designed file consists of coefficients of video pixels in discrete coefficient. The obtained coefficients are transformed to corresponding gray intensity pixel level and coded to binarized coefficient in 16-bit precision. The coefficients are stored in a FIFO device of size 64x16 for burst processing. The FIFO data are packetized based on the coding standard and passed to interfacing unit for bit allocation based on the demanded resource and available resource unit. The incorporation of the packet scheduling interface unit helps in providing scalable coding for multi-bit rate service. The processing unit for the video processing is referred from the standard MPEG structure and a simplified model is realized as shown below,



Figure 2: Digital modelling for coding structure of video processing

The designed system reads the video sequence and interface to the developed digital system via the test interface unit called test bench. The video sequences are read using the analog interface unit and is designed in Matlab tool. The developed function reads the passed video file and transform into successive frame sequence. These frame sequence are read and each coefficient of the read frame are transformed from its gray intensity to equivalent binary sequence. These binary sequences are passed as a input to the designed HDL unit via test bench declaration. The two-frame sequence are loaded as a test vector in this test bench and passed to the HDL for further computation. The first interfacing unit in the Designed Digital system is the noise estimation unit, which is used for the elimination of additional noise coefficients added due to transformation or transmission. The noise estimation unit computes the mean error difference between the two mapping blocks from frame 1 and frame 2. The two blocks are successively compared for distance between the two corresponding pixels and accumulated for the mean error calculation. The error values are calculated for all the blocks in a sequence and the least error value is computed for these obtained error values. This error value is taken as the threshold least mean error value for the noise elimination in motion estimation block. The Least mean error value for the estimation is given by,

 $LMSE_{er} = [mean(least error)^2 + 3 * \sigma(least error)] (1)$ 

For the estimation of motion vector BMA algorithm is realized in full search mode. The operation flow chart for this algorithm is as outlined below,



The following flowchart shows the order of activities carried out at the receiver block.



Figure 3 (b): Flowchart for Receiver block

For the realization of the suggested system the above stated algorithm is designed in digital operational units. The design consist of a comparator unit which generates the absolute difference of the two pixels and compares with the previous error value for the selectivity and accumulate the error value to the current error value in case of minimum of error evaluation. Figure below shows the digital unit realized for the evaluation of this operation.



Figure 4: Digital model of matching algorithm for SAD calculation

The comparison operation is realized by the subtraction of the two input coefficients. The obtained difference value is added and this process is repeated for all the pixel values. On every SAD calculation the previous SAD is compared with the current SAD value for its comparison. From the obtained SAD value the motion vectors are predicted. The SAD is defined by,

 $SAD(d_{1,}d_{2}) = \sum_{(n_{1},n_{2})\in B} |S_{1}(n_{1}, n_{2}, k) - S_{2}(n_{1} + d_{1}, n_{2} + d_{2}), k + 1|$ (2) Where s<sub>1</sub> (n<sub>1</sub>, n<sub>2</sub>, k) is the pixel value at (n<sub>1</sub>, n<sub>2</sub>) in frame k and s<sub>2</sub> (n<sub>1</sub>+d<sub>1</sub>, n<sub>2</sub>+d<sub>2</sub>, k+1) is the pixel value at (n<sub>1</sub>+d<sub>1</sub>, n<sub>2</sub>+d<sub>2</sub>) in frame k+1, d<sub>1</sub> and d<sub>2</sub> are the horizontal and vertical motion vectors respectively. Once the SAD value for a particular frame is calculated the motion vector (d<sub>1</sub>, d<sub>2</sub>) is calculated given by,  $[d_{1}d_{2}]^{T} = arg \min_{d_{1}d_{2}} SAD(d_{1,}d_{2})$ (3)

The developed BMA architecture is interfaced with a general video interface unit providing error correction coding and entropy encoding for its evaluation. The obtained motion vector is passed to entropy encoder and the compressed data is processed with a CRC unit for error coding in the transmitting bit.



Once the compressed bit sequence is obtained and the FCS bits are generated the data is aligned in a frame format before transmission. A general format of the data transmission is developed for the transmission of this data.



Figure 6: Frame format for developed codec

The frame is developed with the following fields of entry, for the generated frame sequence resource allocation is performed. The designing of allocation logic is developed based scheduler logic under multiple level service compatibility. The approach of scalability approach for developed system is outlined in following section.

#### III. Scalable Coding

To provide scalability under different demanded data rate a scheduler scheme is proposed. For the processing of data into end to end streaming a cross layer design is proposed. This system repetitively records the average bandwidth to the streaming server according to the residual BS capacity and RF condition on the period of Report Period. In each of the Report Period, the unit requests a target bit-rate from the coding unit. The data at lower spatial-SNR resolution are given more importance and in each spatial-SNR layer, the lower temporal layer is more important. Therefore, according to the requested bit-rate from the coding unit, the lower spatial-SNR layers are firstly extracted.



Figure 7: video streaming encoder on level selection

The coding structure for such a unit is as shown below to support the data rate at different demanded service rate the coding unit keep a track of the available bandwidth and allocates the resource based on request priority. This results in the coding of smoother data streaming under multiple data rate demand. To achieve the transmission rate faster and retraining the visual quality a projection modelling using linear interpolation is proposed.

# IV. Visual Projection Coding

The projection approach developed codes the resolution under horizontal and vertical grid projection to project the sample the developed projection algorithm is as outlined below. The Pv and PH values defined in equation below.

$$\begin{split} & P_V = \begin{bmatrix} Q_V(X_V/L_V - Z_V) + \varepsilon \end{bmatrix} & (4) \\ & P_H = \begin{bmatrix} Q_H(X_H/L_H - Z_H) + \varepsilon \end{bmatrix} & (5) \\ & \text{Where,} \\ & P_V = \text{Vertical phase.} \\ & P_H = \text{Horizontal phase.} \\ & X_V = \text{Vertical coordinates of high resolution pixel.} \\ & X_H = \text{Horizontal coordinates of high resolution pixel.} \\ & \varepsilon = \text{Very small number like 10-6.} \\ & Z_V = \text{Vertical coordinates of low resolution pixel.} \\ & Z_H = \text{Horizontal coordinates of low resolution pixel.} \\ & For scaling ratios less than two, selects the pixels to be omitted using equation shown below.} \\ & Z_V = \begin{bmatrix} X_V/L_V \end{bmatrix} Z_H = \begin{bmatrix} X_H/L_H \end{bmatrix} & (6) \end{split}$$

The processing data are passed via a input interface unit and the Input memory (IM) unit operates at input pixel clock frequency Fclk\_in. The block consists of four cache buffers, to provide a 5 x 1 pixel column to the CU.

The length of the line buffers is equal to the input video's horizontal size. Shown below is the block diagram of the Input Unit.



#### (a)Operation of Segregator

Segregator unit operates at core clock frequency CU, and generates four high resolution pixels using the selected Fclk\_core. Figure below shows the architecture of the feature extraction and context classification units. Parallel implementation where Dr = 1, will use 15 adders, 32 multipliers, and a serial-parallel implementation, where Dr=4 reduces the number of adders/subtractors to four and the multipliers to eight, with a negligible increase in the number of pipeline registers.



Figure 9: Context Classification and Feature Extraction.

#### (b) Operation of Interpolation

Interpolation unit (IN) unit operates at core clock frequency. 100 multiplications, and 96 additions required to perform 5 x 5 convolution for four HD pixels is resource shared with Dr=4 to reduce the number of multipliers to 25, adders to 24. Therefore at each Fclk\_core clock cycle, one convolution operation is performed, generating four HD pixels at every Fclk-in, clock cycle.

#### (c) Output Unit

Output memory (OM) unit also operates at core clock frequency. It basically chooses the appropriate pixels between the high resolutions pixels generated by IN unit, and arranges them in raster scan order according to the scaling ratio.

#### V. Experimental Results

For the functional evaluation of the designed system the developed system is designed using VHDL definition and simulated on the Active-HDL simulator for its operational verification. For the testing of the designed system a test bench file is generated where two frame data is passed. This frame data is read by the video\_codec file and divide into 8 X 8 Block. The block data is processed for noise estimation, motion estimation, and compression. Under recovery the data is decoded back for its regeneration. The obtained simulation results are as illustrated below. For the real time Realization and resource utilization the developed design is synthesized using Xilinx synthesizer. The obtained synthesis results and their logical, mechanical, and timing reports are also presented in the following section.



Figure 10: simulation result showing the timing simulation for the block division



Figure 11: simulation result illustrating the moving vectors and noise estimation



Figure 12: simulation result illustrating the generated code sequence for the coefficient compression



Figure 13: simulation result illustrating the generation of serial output information

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Figure 14: simulation resulting illustrating the generation of decoded information

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Figure 15: simulation results illustrating buffered code words received and their corresponding decoded word bits.



Figure 16: simulation result illustrating the total generation of the transmitted code sequence and it's decoded bit stream.

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Figure 17: simulation result illustrating the generated frame value after decoding operation

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Figure 18: Illustrating the comparison of the transmitted and recovered frame sequence

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**Figure 19:** simulation result illustrating the CRC operation for the developed system

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**Figure 20:** simulation of frame generation unit

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Figure 21: simulation result for frame reader unit.

Minimum period	: 7.693ns
(Maximum Frequency:	129.984MHz)
Design Statistics	
# IOs	: 34
Cell Usage:	
# BELS	: 542

	<b>Device Utilization</b>	Summary		
Logic Utilization	Used	Available	Utilization	Note(s)
Number of Slice Flip Flops	113	88,192	1%	
Number of 4 input LUTs	123	88,192	1%	
Logic Distribution				
Number of occupied Slices	134	44,096	1%	
Number of Slices containing only related logic	134	134	100%	
Number of Slices containing unrelated logic	0	134	0%	
Total Number of 4 input LUTs	261	88,192	1%	
Number used as logic	123			
Number used as a route-thru	138			
Number of bonded IOBs	34	1,164	2%	
IOB Flip Flops	14			
Number of PPC405s	0	2	0%	
Number of GCLKs	1	16	6%	
Number of GTs	0	20	0%	
Number of GT10s	0	0	0%	
Total equivalent gate count for design	2,633			

Figure 22: summarized synthesis report for the developed desig

Release 9.1i - XPower SoftwareVersion:J.30 Copyright (c) 1995-2007 Xilinx, Inc. All rights reserved. Design: C:\Xilinx91i\BMA\video\_codec.ncd Preferences: video\_codec.pcf Part: 2vpx70ff1704-6 Data version: PREVIEW,v1.0,05-28-03

Power summary:	I(mA)	P(mW)
Total estimated power consumption:		181
Vccint 1.50V :	85	128
Vccaux 2.50V :	20	50
Vcco25 2.50V :	2	4
Clocks :	0	0
Inputs :	0	0
Logic :	0	0
Outputs :		
Vcc :	25	0
Signals :	0	0
Quiescent Vccint 1.50V :	85	128
Quiescent Vccaux 2.50V :	20	50
Quiescent Vcco25 2.50V :	2	4
Thermal summary:		
Estimated junction temperature :	25C	
Ambient temp :	25C	
Case temp :	25C	
Theta J-A :	0C/W	



Figure 23: RTL view of the implemented system using Xilinx synthesizer.

The system is evaluated for (1) Peak Signal-to-Noise Ratio (PSNR) calculated between the predicted motion-compensated frame and the original frame (2) Mean Square Error (MSE) between the predicted frame and the original frame (3) Number of absolute operations per frame for calculating the matching error (4) Execution time according to our implementation and hardware environment.

The obtained observation for certain cases are illustrated below,

Case 1:

Name of the sample: Suzie.qcif Number of frames : 150



Figure 24(a) A frame from Suzie.qcif



Table 1: Performance comparison of FS and DS for suzie image sequence

Algorithm	MSE	PSNR (dB)	Computa	tions	Exe	cution tir	ne(sec)
			number	%	T <sub>1</sub>	$T_2$	$T=T_1+T_2$
FS	15.65	36.18	18271	100	162.13	0.42	162.55
DS	15.91	36.19	1229	6.72	15.11	0.45	15.56

Case 2:

Name of the sample Number of frames : Table tennis.qcif : 150



Figure 25(a): A frame from tabletennis.qcif





Algorithm	MSE	PSNR (Æ)	Computations		Execution ti	me(sec)	
			Number	%	I	Tz	T=T1+T2
FS	85	28.8	18271	100	160.6	0.32	160.92
DS	92	28.4	1303	7.13	16.8	0.32	17.12

Table 2: Performance comparison of FS and DS for table tennis image sequence

# Case 3:

Name of the Sample: Nba.cif Number of frames: 150



Figure 26(a) Nba







Figure 26(c): Frame No. Vs Computations



Figure 26(d): Frame No. Vs PSNR

**Table 3:** Performance Comparison of FS and DS with Nba image sequence

		PSN	Comput	Computations		Execution time(sec)		
Algorithm	MSE	R	Numb	%	T1	T <sub>2</sub>	$T=T_1+T$	
		(dB)	er				2	
FS	522	20.9	80896	100	696.	11.1	707.5	
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DS	612	20.2	8635	10.6	109.	13.4	122.8	
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#### VI. Conclusion

A digital modelling for the processing of video coder under variable resource allocation and demanded data rate is proposed. The developed design provides a simplified approach to the modelling of video processor for high rate processing and higher visual quality observation. The objective of developing a simplified approach of video coder is achieved. The obtained results illustrate the significance of utilizing the proposed system for quality of service in video codec. A relative accuracy in the obtained estimation for video sample at full search is observed. The PSNR value is obtained in improvement to the existing system. The computation time for such coding is observed to be minimized and hence the operational speed is improved in hardware system.

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